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UNITED STATES PATENT AND TRADEMARK OFFICE

APPLICATION NO.	FILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.
09/692,709	10/19/2000	Christian Gehrmann	45687-00036	7545

38065	7590	03/29/2006

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EXAMINER
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ART UNIT	PAPER NUMBER
2136	

DATE MAILED: 03/29/2006

Please find below and/or attached an Office communication concerning this application or proceeding.

-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --

Period for Reply

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE _3_ MONTH(S) OR THIRTY (30) DAYS, WHICHEVER IS LONGER, FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133). Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

Status

1)☒ Responsive to communication(s) filed on _06 January 2006_.

2a)☒ This action is **FINAL.**   2b)☐ This action is non-final.

3)☐ Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

Disposition of Claims

4)☒ Claim(s) _1 and 4-23_ is/are pending in the application.

  4a) Of the above claim(s) _____ is/are withdrawn from consideration.

5)☐ Claim(s) _____ is/are allowed.

6)☒ Claim(s) _1 and 4-23_ is/are rejected.

7)☐ Claim(s) _____ is/are objected to.

8)☐ Claim(s) _____ are subject to restriction and/or election requirement.

Application Papers

9)☐ The specification is objected to by the Examiner.

10)☐ The drawing(s) filed on _____ is/are: a)☐ accepted or b)☐ objected to by the Examiner.

  Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).

  Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).

11)☐ The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

Priority under 35 U.S.C. § 119

12)☐ Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).

  a)☐ All   b)☐ Some * c)☐ None of:

  1.☐ Certified copies of the priority documents have been received.

  2.☐ Certified copies of the priority documents have been received in Application No. _____.

  3.☐ Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).

  * See the attached detailed Office action for a list of the certified copies not received.

Attachment(s)

1) ☐ Notice of References Cited (PTO-892)

2) ☐ Notice of Draftsperson's Patent Drawing Review (PTO-948)

3) ☐ Information Disclosure Statement(s) (PTO-1449 or PTO/SB/08) Paper No(s)/Mail Date _____.

4) ☐ Interview Summary (PTO-413) Paper No(s)/Mail Date. _____ .

5) ☐ Notice of Informal Patent Application (PTO-152)

6) ☐ Other: _____.

DETAILED ACTION

1.   Claims 1 and 4-23 are pending in this office action.

2.   Applicant's arguments, filed January 6, 2006, have been fully considered but they

are not persuasive.

Rejections

3.   The text of those sections of Title 35, U.S. Code not included in this action can

be found in a prior Office action.

Claim Rejections - 35 USC § 103

4.   <u>Claims 1 and 4-6, 17, and 18</u> are rejected under 35 U.S.C. 103(a) as being

unpatentable over <u>Merging and Extending the PGP and PEM Trust Models – The ICE-

TEL Trust Model, Chadwick et al., May/June 1997</u> (hereinafter referred to as Chadwick

et al.) in view of <u>Hunt et al.</u> (U.S. Patent No. 5,539,881).

Regarding <u>claims 1, 4, 5, and 17</u>, <u>Chadwick et al.</u> teaches a method/ad hoc

communication network for establishing security in an ad hoc communication network,

the ad hoc communication network comprising:

- A set of communication nodes (fig. 2, pg. 20),

- At least two nodes of the set of communication nodes having a mutual trust

   relation and comprising a trust group (pg. 20, right column, first full paragraph),

- o The trust relations being created with public keys, and at least one

  additional node (fig. 4, pg. 22),

  - The at least one additional node being a candidate node for joining

    the trust group within the ad hoc communication network (pg. 22,

    left column, first paragraph),

- The nodes having authority to delegate trust to nodes of the set of

  communication nodes within the trust group (pg. 20, "Certification Path"),

- The method comprising the steps of:

  - o Receiving a request from the candidate node to join the trust group within

    said ad hoc communication network wherein said ad hoc communication

    network does not include a separate certificate authority (pg. 20, left

    column, last paragraph. Chadwick teaches (page 19, "Trusted Point") that

    a security domain can be as small as a single user with a user as its

    trusted point, not a certificate authority.); and

  - o Identifying **any node** within the trust group having a trust relation with the

    candidate node (pg. 20, right column, "Cross Certification"), the node

    having the trust relation with the candidate node being an X-node, and

  - o Wherein X-node further sending a signed message comprising a list of

    nodes that the X-node trusts within the ad hoc communication network

    and all corresponding public keys to the candidate node (pg. 20, left

    column, first paragraph).

Chadwick et al. does not teach distributing trust relations between all members in the trust group and the candidate node by means of the X-node distributing the public key associated with said candidate node to said all members of the trust group.

Hunt et al. teaches distributing the public keys and trust relations between all members in the trust group and the candidate node by means of the X-node distributing the public key associated with said candidate node to said all members of the trust group (col. 1, line 42 through col. 2, lines 3).

It would have been obvious to one of ordinary skill in the art, at the time the invention was made, to combine the X-node distributing the public key of the candidate node to all members of the trusted group, as taught by Hunt et al., with the method/network of Chadwick et al. It would have been obvious for such modifications because in the public-key cryptographic system of Chadwick et al., distributing your public key to all the devices/users that you want to be able to communicate with securely, allows the devices/users to encrypt communications with your public key so that only you can decrypt the communications with your secretly held private key. This is desirable because in a wireless ad hoc network, where devices have no base stations or agreed upon topology, a device has the ability to broadcast/publish its public key for any device/user willing to communicate securely with the device. Also, exchanging public keys of every node in the system, communication with other network elements can be accomplished (see col. 1, line 67 through col. 2, line 3 of Hunt et al.).

Regarding claims 6 and 18, Chadwick et al. as modified by Hunt et al. teaches wherein the ad hoc communication network comprises a set of nodes comprising several trust groups (see fig. 4, pg. 22 of Chadwick et al.), each of the set of nodes being candidates for joining all trust groups within the ad hoc communication network that the set of nodes are not already a member of (see pg. 22, left column, second paragraph of Chadwick et al.), the method comprising, after receiving the messages, each node of the set of nodes creating a list of candidate nodes that a given node of the set of nodes trusts and corresponding public keys (see pg. 22, left column, first paragraph of Chadwick).

Claims 7-16 and 19-23 are rejected under 35 U.S.C. 103(a) as being unpatentable over Chadwick et al. in view of Hunt et al. (USPN '881), and further in view of Morris et al. (U.S. Patent No. 6,691,173).

Regarding claims 7 and 19, Chadwick et al./Hunt et al. teaches all the limitations of claims 1, 6 and 17, respectively, above. However, Chadwick et al./Hunt et al. does not teach deciding one node within the ad hoc communication network to act as a server node.

Morris et al. teaches further comprising deciding one node within the ad hoc communication network to act as a server node (col. 4, lines 49-56).

It would have been obvious to one of ordinary skill in the art, at the time the

invention was made, to combine deciding one node to act as a server node, as taught

by Morris et al., with the method/network of Chadwick et al./Hunt et al. It would have

been obvious for such modifications because an ad hoc network needs to establish one

node as the server, while the other nodes act as slaves. This step is necessary and is

therefore a desirable and obvious step.

Regarding claim 8, the combination of Chadwick et al. in view of Hunt et

al./Morris et al. teaches further comprising the server node receiving, from each other

node within the ad hoc communication network, a message comprising a respective

public key, a respective list of candidate nodes that the respective node trusts, and

corresponding public keys (see col. 3, line 49 through col. 4, line 2 of Morris et al.).

Regarding claims 9 and 20, the combination of Chadwick et al. in view of Hunt et

al./Morris et al. teaches further comprising the server node classifying the at least one

candidate node as being a server-trusted node or as being a server-untrusted node,

depending on whether the server node trusts the at least one candidate node or not

(see pg. 22, left column, second paragraph of Chadwick et al.).

Regarding claims 10 and 21, the combination of Chadwick et al. in view of Hunt

et al./Morris et al. teaches wherein the identifying step further comprises the server

node identifying at least one Y-node required for distributing trust relations between the

server node and at least one server-untrusted node (see col. 8, lines 23-37 of Morris et al.).

Regarding claims 11 and 22, the combination of Chadwick et al. in view of Hunt et al./Morris et al. teaches wherein said distributing step further comprises sending, by the server node, of a request to the identified at least one Y-node to distribute said trust relations between the server node and the server-untrusted nodes (see col. 8, lines 38-45 of Morris et al.).

Regarding claim 12, the combination of Chadwick et al. in view of Hunt et al./Morris et al. teaches wherein said distributing step further comprises obtaining, by the server node, of said requested trust relations (see col. 8, lines 45-49 of Morris et al.).

Regarding claim 13, the combination of Chadwick et al. in view of Hunt et al./Morris et al. teaches wherein the step of obtaining the trust relations further comprises:

- Signing, by the Y-node, of the public key of the server node for each server-untrusted node that the Y-node has a trust relation with (see pg. 20, "Certification Path" of Chadwick et al.); and

- Forwarding, by the Y-node, of said signed public key to the server-untrusted node (see pg. 20, "Certification Path" of Chadwick et al.).

Regarding <u>claim 14</u>, the combination of <u>Chadwick et al.</u> in view of <u>Hunt et al.</u>/<u>Morris et al.</u> teaches wherein the step of obtaining the trust relations comprises:

- Signing, by the Y-node, of the public key of the server-untrusted node for each server-untrusted node that the Y-node has a trust relation with (see pg. 20, "Certification Path" of Chadwick et al.); and

- Forwarding, by the Y-node, of said signed public key to the server node (see pg. 20, "Certification Path" of Chadwick et al.).

Regarding <u>claim 15</u>, the combination of <u>Chadwick et al.</u> in view of <u>Hunt et al.</u>/<u>Morris et al.</u> teaches comprising the further step of, after obtaining said trust relation, reclassifying, by the server node, the server-untrusted node with the obtained trust relation as being a server-trusted node (see col. 8, lines 45-49 of Morris et al.).

Regarding <u>claims 16 and 23</u>, the combination of <u>Chadwick et al.</u> in view of <u>Hunt et al.</u>/<u>Morris et al.</u> teaches comprising the further step of sending, by the server node, of a signed message comprising the server node's trusted public keys belonging to trusted candidate nodes within the ad hoc communication network (see col. 4, lines 3-8 of Morris et al. and pg. 20, right column, "Cross Certification" of Chadwick et al.).

Response to Arguments

5.   Applicant argues Hunt et al. does not teach "any node within the trust group having a trust relation with the candidate node being the X-node" and distributing trust

relationships between all members in the trust group and the candidate node by means

of the X-node distributing the public key associated with said candidate node to said all

members of the trust group (page 7, last paragraph through page 8, first paragraph).

Regarding applicant's argument, examiner disagrees. First, Hunt et al. teaches,

on column 3, lines 16-20 that network element 100 may be either a DSNE or a remote

NE. This means that any network element has the opportunity to become the DSNE, or

X-node. Column 3, lines 40-54, shows how the network element behaves if it is the

DSNE or the remote NE. Second, Hunt et al. teaches distributing the information of a

newly added network element to all other nodes in the trust group by the teachings of

figure 5 and column 6, lines 45-58. The cited passage and accompanying figure show

that once a single remote NE (candidate node) is registered with the DSNE (X-node),

the two nodes swap information; then the remote NE (candidate node) receives

information from other nodes and transmits its information to other nodes.

Conclusion

6.   **THIS ACTION IS MADE FINAL.** Applicant is reminded of the extension of time

policy as set forth in 37 CFR 1.136(a).

A shortened statutory period for reply to this final action is set to expire THREE

MONTHS from the mailing date of this action. In the event a first reply is filed within

TWO MONTHS of the mailing date of this final action and the advisory action is not

mailed until after the end of the THREE-MONTH shortened statutory period, then the

shortened statutory period will expire on the date the advisory action is mailed, and any

extension fee pursuant to 37 CFR 1.136(a) will be calculated from the mailing date of

the advisory action. In no event, however, will the statutory period for reply expire later

than SIX MONTHS from the mailing date of this final action.

Any inquiry concerning this communication or earlier communications from the

examiner should be directed to Brandon S. Hoffman whose telephone number is 571-

272-3863. The examiner can normally be reached on M-F 8:30 - 5:00.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's

supervisor, Ayaz R. Sheikh can be reached on 571-272-3795. The fax phone number

for the organization where this application or proceeding is assigned is 571-273-8300.

Information regarding the status of an application may be obtained from the

Patent Application Information Retrieval (PAIR) system. Status information for

published applications may be obtained from either Private PAIR or Public PAIR.

Status information for unpublished applications is available through Private PAIR only.

For more information about the PAIR system, see http://pair-direct.uspto.gov. Should

you have questions on access to the Private PAIR system, contact the Electronic

Business Center (EBC) at 866-217-9197 (toll-free).

CHRISTOPHER REVAK
PRIMARY EXAMINER

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